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# Factor-critical property in 3-dominating-critical graphs ${ }^{\text {T }}$ 

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#### Abstract

A vertex subset $S$ of a graph $G$ is a dominating set if every vertex of $G$ either belongs to $S$ or is adjacent to a vertex of $S$. The cardinality of a smallest dominating set is called the dominating number of $G$ and is denoted by $\gamma(G)$. A graph $G$ is said to be $\gamma$-vertex-critical if $\gamma(G-v)<\gamma(G)$, for every vertex $v$ in $G$.

Let $G$ be a 2-connected $K_{1,5}$-free 3 -vertex-critical graph of odd order. For any vertex $v \in V(G)$, we show that $G-v$ has a perfect matching (except two graphs), which solves a conjecture posed by Ananchuen and Plummer [N. Ananchuen, M.D. Plummer, Matchings in 3-vertex critical graphs: The odd case, Discrete Math., 307 (2007) 1651-1658].


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## 1. Introduction

Let $G$ be a finite simple graph with vertex set $V(G)$ and edge set $E(G)$. A vertex subset $S$ of $G$ is a dominating set if every vertex of $G$ either belongs to $S$ or is adjacent to a vertex of $S$. The minimum size of such a set is called the dominating number of $G$ and is denoted by $\gamma(G)$. A graph $G$ is vertex-domination-critical, or $\gamma$-vertex-critical, if for any vertex $v$ of $G, \gamma(G-v)<\gamma(G)$. We use $G[S]$ to denote the subgraph induced by $S$ for some $S \subseteq V(G)$. The minimum degree of $G$ is denoted by $\delta(G)$. A graph is called $K_{1, k}-$ free if it has no induced subgraph isomorphic to the complete bipartite graph $K_{1, k}$.

A set of independent edges in a graph $G$ is called a matching of $G$. A matching is perfect if it is incident with every vertex of $G$. If $G-v$ has a perfect matching, for every choice of $v \in V(G), G$ is said to be factor-critical. The concept of factor-critical graphs was first introduced by Gallai in 1963 and it plays an important role in the study of matching theory. In contrast to its apparent strong property, such graphs form a relatively rich family for study. It is the essential "building block" for the so-called Gallai-Edmonds structure of the graphs with matchings.

The subject of $\gamma$-vertex-critical graphs was studied first by Brigham, Chinn and Dutton [3,4] and continued by Fulman et al. [5,6]. Clearly, the only 1 -vertex-critical graph is $K_{1}$ (a single vertex). Brigham, Chinn and Dutton [3]

[^0]pointed out that the 2-vertex-critical graphs are precisely the family of graphs obtained from the complete graphs $K_{2 n}$ with a perfect matching removed. For $\gamma>2$, however, much remains unknown about the structure of $\gamma$-vertex-critical graphs. Recently, Ananchuen and Plummer [1,2] began to study matchings in 3-vertex-critical graphs. They showed that a $K_{1,5}$-free 3-vertex-critical graph of even order has a perfect matching (see [1]) and a $K_{1,4}$-free 3-vertex-critical graph of odd order is factor-critical (see [2]). Furthermore, they posed the following conjecture.

Conjecture 1. If $G$ is a $K_{1,5}$-free 3-vertex-critical 2-connected graph of odd order with $\delta(G) \geq 3$, then $G$ is factorcritical.

In this paper, we show that the conjecture holds for almost all graphs and there are only two counterexamples.
Let $v \in V(G)$, we denote a minimum dominating set of $G-v$ by $D_{v}$. The following facts about $D_{v}$ follow immediately from the definition of 3-vertex-criticality and we shall use it frequently in the proof of the main theorem.
Facts: If $G$ is 3-vertex-critical, then the following hold
(1) For every vertex $v$ of $G,\left|D_{v}\right|=2$.
(2) If $D_{v}=\{x, y\}$, then $x$ and $y$ are not adjacent to $v$.
(3) For every pair of distinct vertices $v$ and $w, D_{v} \neq D_{w}$.

The readers are referred to [7] for other terminologies not specified in this paper.

## 2. Main result

By Tutte's well-known 1-Factor Theorem, if a graph $G$ has no perfect matching, then there exits a set $S \subseteq V(G)$ such that the number of components in $G-S$ having odd order is greater than the order of $S$. Let $S \subseteq V(G)$, we shall denote by $\omega(G-S)$, the number of components of $G-S$ and by $c_{o}(G-S)$, the number of odd components of $G-S$. A criterion similar to 1-Factor Theorem for factor-critical graphs is as follows.

Lemma 2.1 (see [7]). A graph $G$ is factor-critical if and only if $c_{o}(G-S) \leq|S|-1$, for every nonempty set $S \subseteq V(G)$.

Lemma 2.2. Let $G$ be 3-vertex-critical and $S$ be a cutset in $G$ with $|S| \geq 4$. If $D_{u} \subseteq S$ for each vertex $u \in S$, then there exists no vertex of degree 1 in $G[S]$.

Proof. Suppose to the contrary that there exists some $v \in S$ such that $v$ is of degree 1 in $G[S]$. Without loss of generality, let $v w \in E(G)$, where $w \in S$. By Fact 2 , $v \notin D_{w}$. Since $D_{w} \subseteq S, D_{w}$ does not dominate $v$, a contradiction.

The following two lemmas, proved by Ananchuen and Plummer [2], will be used in our proof of the main theorem.
Lemma 2.3. If $G$ is 3 -vertex-critical and $S$ is a cutset in $G$ such that $\omega(G-S) \geq 4$ or $\omega(G-S)=3$, but each component has at least 2 vertices, then each vertex of $G-S$ is not adjacent to at least one vertex of $S$.

Lemma 2.4. Let $G$ be a 3-vertex-critical graph and suppose that $S$ is a cutset of order 2 in $G$, then $\omega(G-S) \leq 3$. Furthermore, if $\omega(G-S)=3$, then $G-S$ must contain at least one singleton component.

Before giving our main result, we note that the graphs $G_{1}$ and $G_{2}$ in Fig. 1 are $K_{1,5}$-free 3-vertex-critical 2connected graph of order 11 with $\delta(G)=3$, but are not factor-critical, since $G_{i}-v_{i}$ has no perfect matching for $i=1,2$. We shall show that these two graphs are the only two counterexamples to Conjecture 1 .

Theorem 2.1. If $G$ is a $K_{1,5-f r e e ~ 3-v e r t e x-c r i t i c a l ~ 2-c o n n e c t e d ~ g r a p h ~ o f ~ o d d ~ o r d e r ~ w i t h ~} \delta(G) \geq 3$, except the graphs $G_{1}$ and $G_{2}$ shown in Fig. 1, then $G$ is factor-critical.

Proof. Suppose that $G$ is not factor-critical. By Lemma 2.1 and the parity, there exists a nonempty set $S \subseteq V(G)$ such that $c_{o}(G-S) \geq|S|+1$. Without loss of generality, let $S$ be a minimal such set with $|S|=k$. Then $k \geq 2$ as $G$ is 2 -connected. Let $C_{1}, C_{2}, \ldots, C_{t}$ be the odd components of $G-S$ and $E_{1}, E_{2}, \ldots, E_{n}$ the even components of $G-S$. We consider the following cases.


Fig. 1. The graphs $G_{1}$ and $G_{2}$.
Case 1. $k=2$.
By Lemma 2.4, then $t=3$ and $G-S$ has no even components. Since $\delta(G) \geq 3$ and $k=2$, each odd component of $G-S$ has at least three vertices, which contradicts Lemma 2.4.

Case 2. $k=3$.
Thus, $t \geq 4$. By Lemma 2.3, each vertex of $G-S$ is not adjacent to at least one vertex of $S$. Since $\delta(G) \geq 3$ and $k=3$, we have $\left|V\left(C_{i}\right)\right| \geq 3$ for $i=1,2, \ldots, t$. By Fact 3, there must exist a vertex $x$ in some odd component of $G-S$ such that $D_{x} \nsubseteq S$. Clearly, $D_{x} \cap S \neq \emptyset$. Without loss of generality, let $x \in V\left(C_{1}\right)$ and $D_{x}=\{u, y\}$, where $u \in S$ and $y \in V(G)-S$. Since $G$ is $K_{1,5}$-free, by the parity, so $t=4$ and $G-S$ has at most one even component.

Claim 1. There exists an odd component $C_{j}(j \geq 2)$ such that $C_{j}$ is a complete graph and $u$ dominates $V\left(C_{j}\right)$.
If $y \in V\left(C_{1}\right)-\{x\}$, then $u$ is adjacent to every vertex of $\bigcup_{i=2}^{4} V\left(C_{i}\right)$. Since $G$ is $K_{1,5-f r e e, ~ a t ~ l e a s t ~ t w o ~}^{\text {, }}$ of $C_{2}, C_{3}$ and $C_{4}$ are complete. If $y \in \bigcup_{i=2}^{4} V\left(C_{i}\right)$, say $y \in V\left(C_{2}\right)$, then $u$ dominates all the vertices of $\left(V\left(C_{1}\right) \cup V\left(C_{3}\right) \cup V\left(C_{4}\right)\right)-\{x\}$, and at least one of $C_{3}$ and $C_{4}$ is complete, by $K_{1,5}$-freeness in $G$ again. If $G-S$ has an even component $E_{1}$ and $y \in V\left(E_{1}\right)$, then $u$ is adjacent to every vertex of $\bigcup_{i=1}^{4} V\left(C_{i}\right)-\{x\}$. Since $G$ is $K_{1,5}$-free, $C_{2}, C_{3}$ and $C_{4}$ are all complete. So Claim 1 is proved.

Without loss of generality, assume that $C_{4}$ is complete and $u$ dominates $V\left(C_{4}\right)$.
Claim 2. Each vertex of $S-\{u\}$ is not adjacent to any vertex of $V\left(C_{4}\right)$.
Suppose to the contrary that va$a_{4} \in E(G)$ for some $v \in S-\{u\}$ and $a_{4} \in V\left(C_{4}\right)$. Then $D_{a_{4}} \cap\left(\{u, v\} \cup V\left(C_{4}\right)\right)=\emptyset$, as $C_{4}$ is complete and $u a_{4} \in E(G)$. Let $S-\{u, v\}=\{w\}$. Clearly, $w \in D_{a_{4}}$. Then $w a_{4} \notin E(G)$ and $w$ dominates $V\left(C_{4}\right)-\left\{a_{4}\right\}$. Let $b_{4} \in V\left(C_{4}\right)-\left\{a_{4}\right\}$. Then $u b_{4} \in E(G)$ and $w b_{4} \in E(G)$. Consequently, $D_{b_{4}} \cap\left(\{u, w\} \cup V\left(C_{4}\right)\right)=\emptyset$ and $v \in D_{b_{4}}$. So $v b_{4} \notin E(G)$ and $v$ dominates $V\left(C_{4}\right)-\left\{b_{4}\right\}$. Now let $c_{4} \in V\left(C_{4}\right)-\left\{a_{4}, b_{4}\right\}$, then $c_{4}$ is adjacent to every vertex of $S$, which contradicts Lemma 2.3.

From Claim 2, $u$ is a cutvertex in $G$, which is against the fact that $G$ is 2-connected.
Case 3. $k=4$.
Thus, $t \geq 5$. We first show that there exists some $a \in S$ such that $D_{a} \nsubseteq S$. Otherwise, $D_{b} \subseteq S$ for each vertex $b \in S$. By Lemma 2.2 and Fact 2, every vertex of $S$ in $G[S]$ has degree 0 . It is easy to check that this is impossible.

So let $u \in S$ such that $D_{u} \nsubseteq S$. Clearly, $D_{u} \cap S \neq \emptyset$. Let $D_{u}=\{v, x\}$, where $v \in S$ and $x \in V(G)-S$. Since $G$ is $K_{1,5}$-free, $t=5$ and $G-S$ has no even components. Without loss of generality, let $x \in V\left(C_{1}\right)$, then $v$ dominates all vertices of $\bigcup_{i=2}^{5} V\left(C_{i}\right)$. Moreover, by $K_{1,5}$-freeness again, $C_{2}, C_{3}, C_{4}$ and $C_{5}$ are all complete, and $v$ is not adjacent to any vertex of $V\left(C_{1}\right)$.

Claim 3. Each vertex of $S$ is adjacent to at least three odd components of $G-S$.
Otherwise, there exists a vertex $c \in S$ such that $c$ is adjacent to at most two odd components of $G-S$. Let $S^{\prime}=S-\{c\}$. It is easy to see that $S^{\prime}$ is a nonempty set which satisfies the condition that $c_{o}\left(G-S^{\prime}\right) \geq\left|S^{\prime}\right|+1$, contradicting to the minimality of $S$.

Let $S-\{u, v\}=\{w, z\}$. By Claim 3, wis adjacent to at least two of $C_{2}, C_{3}, C_{4}$ and $C_{5}$. Without loss of generality, let $w c_{i} \in E(G)$, where $c_{i} \in V\left(C_{i}\right)$ for $i=2,3$. Then $z \in D_{c_{2}}$. Otherwise, $u \in D_{c_{2}}$ and $D_{c_{2}} \cap V\left(C_{1}\right) \neq \emptyset$ since $u x \notin E(G)$. But then $D_{c_{2}}$ cannot dominate $v$, a contradiction. Similarly, $z \in D_{c_{3}}$. Thus, $z c_{i} \notin E(G)$ for $i=2$, 3. By Fact 3, then either $D_{c_{2}} \neq\{u, z\}$ or $D_{c_{3}} \neq\{u, z\}$, say $D_{c_{2}} \neq\{u, z\}$. Since $z c_{3} \notin E(G)$, it follows that $D_{c_{2}} \cap V\left(C_{3}\right) \neq \emptyset$ and $z$ dominates every vertex of $V\left(C_{1}\right) \cup V\left(C_{4}\right) \cup V\left(C_{5}\right)$. By similar arguments, $w \in D_{c_{4}}, w \in D_{c_{5}}$
for some $c_{4} \in V\left(C_{4}\right)$ and $c_{5} \in V\left(C_{5}\right)$. Furthermore, $w c_{i} \notin E(G)$ for $i=4,5$, and $w$ is adjacent to all vertices of $V\left(C_{1}\right) \cup V\left(C_{2}\right) \cup V\left(C_{3}\right)$.

We next show that $C_{2}$ is a singleton. Otherwise, $\left|V\left(C_{2}\right)\right| \geq 3$ and let $a_{2}, b_{2} \in V\left(C_{2}\right)-\left\{c_{2}\right\}$. By similar arguments as above, $z \in D_{a_{2}}, z \in D_{b_{2}}$ and either $D_{a_{2}} \neq\{u, z\}$ or $D_{b_{2}} \neq\{u, z\}$. Assume that $D_{a_{2}} \neq\{u, z\}$. Then $D_{a_{2}} \cap V\left(C_{3}\right) \neq \emptyset$, since $z c_{3} \notin E(G)$. But then $z$ is adjacent to all vertices of $V\left(C_{2}\right)-\left\{a_{2}\right\}$ and this contradicts the fact that $z c_{2} \notin E(G)$. Similarly, $C_{3}, C_{4}$ and $C_{5}$ are all singletons of $G-S$. Since $\delta(G) \geq 3, u c_{i} \in E(G)$ for $i=2,3,4,5$. As $G$ is $K_{1,5}$-free, $u$ is not adjacent to any vertex of $V\left(C_{1}\right)$.

Because $\delta(G) \geq 3$ and $u, v$ are not adjacent to any vertex of $V\left(C_{1}\right)$, we have $\left|V\left(C_{1}\right)\right| \geq 3$. Moreover, $D_{x} \cap\left(V\left(C_{1}\right)-\{x\}\right) \neq \emptyset$ and $D_{x} \cap\{u, v\} \neq \emptyset$ (say, $u \in D_{x}$ ). Recall that $u v \notin E(G)$ and $v$ is not adjacent to any vertex of $V\left(C_{1}\right)$, thus $v$ is not dominated by $D_{x}$, a contradiction.

Case 4. $k=5$.
Claim 4. For every vertex $x \in S, D_{x} \subseteq S$.
Otherwise, $D_{u} \nsubseteq S$ for some $u \in S$. Clearly, $D_{u} \cap S \neq \emptyset$. Let $D_{u}=\{y, z\}$, where $y \in S$ and $z \in V(G)-S$. Since $t \geq 6, y$ must dominate at least 5 odd components of $G-S$, which contradicts the fact that $G$ is $K_{1,5}$-free.

Let $S=\left\{s_{1}, s_{2}, s_{3}, s_{4}, s_{5}\right\}$. By Fact 3, there are $\binom{5}{2}=10$ distinct pairs of vertices in $S$ and at least 11 vertices in $G$. So there must exist a vertex $x$ in an odd component of $G-S$ such that $D_{x} \nsubseteq S$. Assume that $x \in V\left(C_{1}\right)$. Clearly, $D_{x} \cap S \neq \emptyset$. Since $G$ is $K_{1,5}$-free, we have $t=6$ and $G-S$ has no even components. By Claim 4 and Lemma 2.2, each vertex of $S$ in $G[S]$ has degree 0 or 2 . It is not hard to see that $G[S]$ can only be a 5 -cycle or a union of a 4-cycle and an isolated vertex.

Case 4.1. $G[S]$ is a 5-cycle.
Let $s_{1} s_{2} s_{3} s_{4} s_{5} s_{1}$ be the 5-cycle in the counterclockwise order and $D_{x}=\left\{s_{1}, w\right\}$, where $w \in V(G)-S$. Since $G$ is $K_{1,5}$-free, $w \notin V\left(C_{1}\right)$. Assume that $w \in V\left(C_{2}\right)$. Then $s_{1}$ is adjacent to all vertices of $\bigcup_{i=3}^{6} V\left(C_{i}\right)$ and $w$ dominates $s_{3}, s_{4}$. Moreover, $K_{1,5}$-freeness of $G$ implies that $C_{3}, C_{4}, C_{5}$ and $C_{6}$ are all complete, $C_{1}$ is a singleton and $s_{1}$ is not adjacent to any vertex of $V\left(C_{1}\right) \cup V\left(C_{2}\right)$.

Since $D_{s_{3}}=\left\{s_{1}, s_{5}\right\}, s_{5}$ is adjacent to each vertex of $V\left(C_{1}\right) \cup V\left(C_{2}\right)$. Similarly, since $D_{s_{4}}=\left\{s_{1}, s_{2}\right\}, s_{2}$ is adjacent to each vertex of $V\left(C_{1}\right) \cup V\left(C_{2}\right)$. Therefore, $w$ is adjacent to all vertices of $S-\left\{s_{1}\right\}$. Now consider $D_{w}$. Since $D_{w} \cap S=\left\{s_{1}\right\}$ and $s_{1} x \notin E(G)$, it follows that $D_{w}=\left\{s_{1}, x\right\}$. Hence, $x$ dominates $s_{3}, s_{4}$ and $V\left(C_{2}\right)=\{w\}$. But then $\left\{s_{1}, s_{3}\right\}$ is a dominating set in $G$, contradicting the assumption that $\gamma(G)=3$.

Case 4.2. $G[S]$ is a union of a 4-cycle and an isolated vertex.
Let $s_{1} s_{2} s_{3} s_{4} s_{1}$ be the 4-cycle in the counterclockwise order and $s_{5}$ the isolated vertex in $G[S]$. Then $D_{s_{1}}=\left\{s_{3}, s_{5}\right\}$, $D_{s_{2}}=\left\{s_{4}, s_{5}\right\}, D_{s_{3}}=\left\{s_{1}, s_{5}\right\}$, and $D_{s_{4}}=\left\{s_{2}, s_{5}\right\}$.

Since $G$ is $K_{1,5}$-free, $s_{5}$ is adjacent to at most 4 odd components of $G-S$. Without loss of generality, let $C_{1}, \ldots, C_{r}$ be the components which are not adjacent to $s_{5}$. Then $t=6$ implies $r \geq 2$. Thus $s_{i}$ is adjacent to every vertex of $\bigcup_{j=1}^{r} V\left(C_{j}\right)$ for $i=1,2,3,4$. Now consider $D_{y}, y \in V\left(C_{1}\right)$. Clearly, $D_{y} \cap S=\left\{s_{5}\right\}$. Since $s_{5}$ cannot dominate $V\left(C_{2}\right), D_{y} \cap V\left(C_{2}\right) \neq \emptyset$. Therefore, $r=2$ and $s_{5}$ is adjacent to every vertex of $\bigcup_{i=3}^{6} V\left(C_{i}\right)$. Moreover, $V\left(C_{1}\right)=\{y\}$. By a similar argument, $C_{2}$ is also a singleton. For each vertex $v \in \bigcup_{i=3}^{6} V\left(C_{i}\right)$, we have $D_{v} \cap S \neq \emptyset$ and $D_{v} \nsubseteq S$, since $s_{5} \notin D_{v}$ and the vertices in $S-\left\{s_{5}\right\}$ do not dominate $s_{5}$. From $K_{1,5}$-freeness of $G, C_{3}, C_{4}, C_{5}$ and $C_{6}$ are all singletons, say $V\left(C_{i}\right)=\left\{c_{i}\right\}$ for $i=3,4,5,6$.

Let $H$ be the induced subgraph in $G$ with vertex set $\left\{s_{i}, c_{j} \mid 1 \leq i \leq 4,3 \leq j \leq 6\right\}$ by deleting the edges in $G[S]$. For $3 \leq j \leq 6$, since $\delta(G) \geq 3, c_{j}$ is adjacent to at least two vertices of $S-\left\{s_{5}\right\}$. On the other hand, since $G$ is $K_{1,5}$-free, each vertex of $S-\left\{s_{5}\right\}$ is adjacent to at most two vertices of $\bigcup_{i=3}^{6}\left\{c_{i}\right\}$. Thus $H$ is a 2-regular bipartite graph and hence consists of either a 8-cycle or a union of two 4-cycles. However, there are only four such graphs under the isomorphism (see Figs. 1 and 2). It is easy to see that $G_{3}$ and $G_{4}$ are not 3-vertex-critical, since $\left|\gamma\left(G_{i}-v_{i}\right)\right|>2$ in $G_{i}$ for $i=3,4$. Therefore, $G_{1}$ and $G_{2}$ are two counterexamples to Conjecture 1.

Case 5. $k \geq 6$.
Claim 5. For every vertex $x \in V(G), D_{x} \subseteq S$.
Suppose that $D_{x} \nsubseteq S$ for some $x \in V(G)$. Clearly, $D_{x} \cap S \neq \emptyset$. Let $D_{x}=\{y, z\}$, where $y \in S$ and $z \in V(G)-S$. Since $t \geq 7$, y must dominate at least 5 odd components of $G-S$, a contradiction to $K_{1,5}$-freeness.


Fig. 2. The graphs $G_{3}$ and $G_{4}$.
Let $w$ be any vertex in $S$, then $D_{w} \subseteq S$ by Claim 5. Since $G$ is $K_{1,5}$-free, each vertex of $D_{w}$ can dominate at most 4 components of $G-S$, which implies that the number of components of $G-S$ is at most 8 or $t \leq 8$. That is, $6 \leq k \leq 7$.

Let $S_{i} \subseteq S$ be the set of vertices in $S$ which are adjacent to some vertex in $C_{i}$ for $i=1,2, \ldots, t$, and let $d=\min \left\{\left|S_{i}\right|\right\}$. Without loss of generality, assume that $\left|S_{1}\right|=d$. Note that for any vertex $v \in V(G)-V\left(C_{1}\right)$, $D_{v} \cap S_{1} \neq \emptyset$. We call such a set $D_{v}$ normal 2-set associated with $v$ and $S_{1}$, or normal set in short. By a simple counting, we see that there are at most $\binom{k}{2}-\binom{k-d}{2}$ normal sets. Since $\left|V(G)-V\left(C_{1}\right)\right| \geq 2 k$, Fact 3 implies $\binom{k}{2}-\binom{k-d}{2} \geq 2 k$ or $d \geq 3$. On the other hand, since $G$ is $K_{1,5}$-free, each vertex of $S$ is adjacent to at most 4 components of $G-S$, that is, $d \leq \frac{4 k}{k+1}$ or $d \leq 3$. Hence $d=3$.

Case 5.1. $k=6$.
Thus $t=7$ and $G-S$ has at most one even component. By Claim 5, there are $\binom{6}{2}=15$ distinct pairs of vertices in $S$ and at least 13 vertices in $G$. So by Fact $3,|V(G)|=13$ or 15 , and $G-S$ has at least 6 singletons.

It is not hard to see that there exists at least four odd components whose corresponding $S_{i}$ 's having the order exactly 3 , and at least two of them are singletons. Without loss of generality, let $C_{1}=\left\{c_{1}\right\}$ and $C_{2}=\left\{c_{2}\right\}$ be two singletons. Then, for every vertex $v \in V(G)-\left\{c_{1}\right\}, D_{v} \cap S_{1} \neq \emptyset$. There are 12 normal sets associated with $S_{1}$ in $S$, and thus $|V(G)|=13$. Next consider $S_{2}$. If $S_{2}=S_{1}$, then $D_{c_{2}}$ cannot dominate $c_{1}$, a contradiction. If $\left|S_{2} \cap S_{1}\right| \leq 2$, however, there must exist 2 normal sets associated with $S_{1}$ which are not adjacent to $c_{2}$, at most one can be realized as $D_{c_{2}}$, and the other cannot dominate $c_{2}$, again a contradiction.

Case 5.2. $k=7$.
Thus $t=8$ and $G-S$ has no even components. By a similar argument that is used in the proof of Case 5.1 , one reaches the same contradiction.

This completes the proof of our theorem.

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